

Lattices of Elementary Substructures

Athar Abdul-Quader

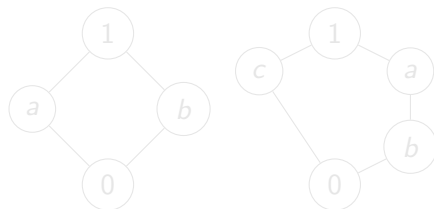
February 19, 2016

Definition

A **lattice** is a partial order $(L, <)$ such that any two elements have a sup (\vee) and an inf (\wedge) .

Example

B_2 : Boolean Algebra on two elements, N_5 : pentagon

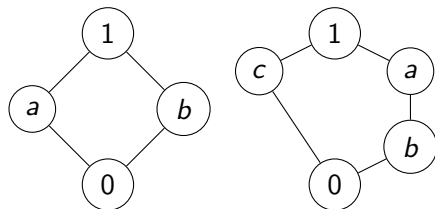


Definition

A **lattice** is a partial order $(L, <)$ such that any two elements have a sup (\vee) and an inf (\wedge) .

Example

\mathbf{B}_2 : Boolean Algebra on two elements, \mathbf{N}_5 : pentagon



Definition

Let L be a lattice.

- ▶ We say $a \in L$ is **compact** if whenever $X \subseteq L$ and $a \leq \bigvee X$, there is a finite $Y \subseteq X$ such that $a \leq \bigvee Y$.
- ▶ L is **algebraic** if it is complete (for all $X \subseteq L$, $\bigvee X$ and $\bigwedge X$ exist), and each element of L is the supremum of a set of compact elements.
- ▶ Let κ be an infinite cardinal. L is **κ -algebraic** if it is algebraic and, whenever $x \in L$ is compact, $|\{a \in L \mid a \leq x \text{ and } a \text{ is compact}\}| < \kappa$.

Definition

Let L be a lattice.

- ▶ We say $a \in L$ is **compact** if whenever $X \subseteq L$ and $a \leq \bigvee X$, there is a finite $Y \subseteq X$ such that $a \leq \bigvee Y$.
- ▶ L is **algebraic** if it is complete (for all $X \subseteq L$, $\bigvee X$ and $\bigwedge X$ exist), and each element of L is the supremum of a set of compact elements.
- ▶ Let κ be an infinite cardinal. L is **κ -algebraic** if it is algebraic and, whenever $x \in L$ is compact, $|\{a \in L \mid a \leq x \text{ and } a \text{ is compact}\}| < \kappa$.

Definition

Let L be a lattice.

- ▶ We say $a \in L$ is **compact** if whenever $X \subseteq L$ and $a \leq \bigvee X$, there is a finite $Y \subseteq X$ such that $a \leq \bigvee Y$.
- ▶ L is **algebraic** if it is complete (for all $X \subseteq L$, $\bigvee X$ and $\bigwedge X$ exist), and each element of L is the supremum of a set of compact elements.
- ▶ Let κ be an infinite cardinal. L is **κ -algebraic** if it is algebraic and, whenever $x \in L$ is compact, $|\{a \in L \mid a \leq x \text{ and } a \text{ is compact}\}| < \kappa$.

Substructure and Interstructure Lattices

Given a model $\mathcal{M} \models \text{PA}$, the set of all $\mathcal{K} \preceq \mathcal{M}$ forms a lattice under the inclusion relation. Additionally, if $\mathcal{M} \prec \mathcal{N}$, then the set of all \mathcal{K} such that $\mathcal{M} \preceq \mathcal{K} \preceq \mathcal{N}$ also forms a lattice.

Definition

Given a model $\mathcal{M} \models \text{PA}$, the set $\{\mathcal{K} \models \text{PA} \mid \mathcal{K} \preceq \mathcal{M}\}$ is called the **substructure lattice** of \mathcal{M} and is denoted $\text{Lt}(\mathcal{M})$. If $\mathcal{M} \prec \mathcal{N}$, then the set $\{\mathcal{K} \mid \mathcal{M} \preceq \mathcal{K} \preceq \mathcal{N}\}$ is called the **interstructure lattice** (between \mathcal{M} and \mathcal{N}) and is denoted $\text{Lt}(\mathcal{N}/\mathcal{M})$.

Example

The simplest examples are any prime model, such as \mathbb{N} . $\text{Lt}(\mathbb{N}) = \mathbf{1}$, the single element lattice.

Substructure and Interstructure Lattices

Given a model $\mathcal{M} \models \text{PA}$, the set of all $\mathcal{K} \preceq \mathcal{M}$ forms a lattice under the inclusion relation. Additionally, if $\mathcal{M} \prec \mathcal{N}$, then the set of all \mathcal{K} such that $\mathcal{M} \preceq \mathcal{K} \preceq \mathcal{N}$ also forms a lattice.

Definition

Given a model $\mathcal{M} \models \text{PA}$, the set $\{\mathcal{K} \models \text{PA} \mid \mathcal{K} \preceq \mathcal{M}\}$ is called the **substructure lattice** of \mathcal{M} and is denoted $\text{Lt}(\mathcal{M})$. If $\mathcal{M} \prec \mathcal{N}$, then the set $\{\mathcal{K} \mid \mathcal{M} \preceq \mathcal{K} \preceq \mathcal{N}\}$ is called the **interstructure lattice** (between \mathcal{M} and \mathcal{N}) and is denoted $\text{Lt}(\mathcal{N}/\mathcal{M})$.

Example

The simplest examples are any prime model, such as \mathbb{N} . $\text{Lt}(\mathbb{N}) = \mathbf{1}$, the single element lattice.

Substructure and Interstructure Lattices

Given a model $\mathcal{M} \models \text{PA}$, the set of all $\mathcal{K} \preceq \mathcal{M}$ forms a lattice under the inclusion relation. Additionally, if $\mathcal{M} \prec \mathcal{N}$, then the set of all \mathcal{K} such that $\mathcal{M} \preceq \mathcal{K} \preceq \mathcal{N}$ also forms a lattice.

Definition

Given a model $\mathcal{M} \models \text{PA}$, the set $\{\mathcal{K} \models \text{PA} \mid \mathcal{K} \preceq \mathcal{M}\}$ is called the **substructure lattice** of \mathcal{M} and is denoted $\text{Lt}(\mathcal{M})$. If $\mathcal{M} \prec \mathcal{N}$, then the set $\{\mathcal{K} \mid \mathcal{M} \preceq \mathcal{K} \preceq \mathcal{N}\}$ is called the **interstructure lattice** (between \mathcal{M} and \mathcal{N}) and is denoted $\text{Lt}(\mathcal{N}/\mathcal{M})$.

Example

The simplest examples are any prime model, such as \mathbb{N} . $\text{Lt}(\mathbb{N}) = \mathbf{1}$, the single element lattice.

Substructure and Interstructure Lattices II

Let $\mathcal{M} \models \text{PA}$. Then the compact elements of $\text{Lt}(\mathcal{M})$ are the finitely generated submodels.

Theorem

$\text{Lt}(\mathcal{M})$ is \aleph_1 -algebraic.

Theorem

If $\mathcal{M} \prec \mathcal{N}$ and $|M| = \kappa$, then $\text{Lt}(\mathcal{N}/\mathcal{M})$ is κ^+ -algebraic.

These are (so far) the only restrictions we currently have for a lattice L to be a substructure lattice or an interstructure lattice.

Substructure and Interstructure Lattices II

Let $\mathcal{M} \models \text{PA}$. Then the compact elements of $\text{Lt}(\mathcal{M})$ are the finitely generated submodels.

Theorem

$\text{Lt}(\mathcal{M})$ is \aleph_1 -algebraic.

Theorem

If $\mathcal{M} \prec \mathcal{N}$ and $|M| = \kappa$, then $\text{Lt}(\mathcal{N}/\mathcal{M})$ is κ^+ -algebraic.

These are (so far) the only restrictions we currently have for a lattice L to be a substructure lattice or an interstructure lattice.

Substructure and Interstructure Lattices II

Let $\mathcal{M} \models \text{PA}$. Then the compact elements of $\text{Lt}(\mathcal{M})$ are the finitely generated submodels.

Theorem

$\text{Lt}(\mathcal{M})$ is \aleph_1 -algebraic.

Theorem

If $\mathcal{M} \prec \mathcal{N}$ and $|M| = \kappa$, then $\text{Lt}(\mathcal{N}/\mathcal{M})$ is κ^+ -algebraic.

These are (so far) the only restrictions we currently have for a lattice L to be a substructure lattice or an interstructure lattice.

Substructure and Interstructure Lattices II

Let $\mathcal{M} \models \text{PA}$. Then the compact elements of $\text{Lt}(\mathcal{M})$ are the finitely generated submodels.

Theorem

$\text{Lt}(\mathcal{M})$ is \aleph_1 -algebraic.

Theorem

If $\mathcal{M} \prec \mathcal{N}$ and $|\mathcal{M}| = \kappa$, then $\text{Lt}(\mathcal{N}/\mathcal{M})$ is κ^+ -algebraic.

These are (so far) the only restrictions we currently have for a lattice L to be a substructure lattice or an interstructure lattice.

- ▶ Gaifman (1960s, published 1976): Every model $\mathcal{M} \models \text{PA}$ has a **minimal** elementary end extension
- ▶ Blass (1974): Every countable nonstandard \mathcal{M} has a minimal cofinal extension.

Definition

A lattice L is called **distributive** if whenever $a, b, c \in L$,
 $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$.

- ▶ Gaifman: Given a finite **distributive** lattice D , there is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong D$.
- ▶ Mills (1979): If D is an \aleph_1 -algebraic distributive lattice, and $\mathcal{M} \models \text{PA}$, there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong D$.

- ▶ Gaifman (1960s, published 1976): Every model $\mathcal{M} \models \text{PA}$ has a **minimal** elementary end extension
- ▶ Blass (1974): Every countable nonstandard \mathcal{M} has a minimal cofinal extension.

Definition

A lattice L is called **distributive** if whenever $a, b, c \in L$,
 $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$.

- ▶ Gaifman: Given a finite **distributive** lattice D , there is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong D$.
- ▶ Mills (1979): If D is an \aleph_1 -algebraic distributive lattice, and $\mathcal{M} \models \text{PA}$, there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong D$.

- ▶ Gaifman (1960s, published 1976): Every model $\mathcal{M} \models \text{PA}$ has a **minimal** elementary end extension
- ▶ Blass (1974): Every countable nonstandard \mathcal{M} has a minimal cofinal extension.

Definition

A lattice L is called **distributive** if whenever $a, b, c \in L$,
 $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$.

- ▶ Gaifman: Given a finite **distributive** lattice D , there is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong D$.
- ▶ Mills (1979): If D is an \aleph_1 -algebraic distributive lattice, and $\mathcal{M} \models \text{PA}$, there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong D$.

- ▶ Gaifman (1960s, published 1976): Every model $\mathcal{M} \models \text{PA}$ has a **minimal** elementary end extension
- ▶ Blass (1974): Every countable nonstandard \mathcal{M} has a minimal cofinal extension.

Definition

A lattice L is called **distributive** if whenever $a, b, c \in L$,
 $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$.

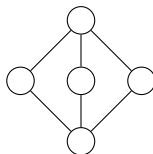
- ▶ Gaifman: Given a finite **distributive** lattice D , there is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong D$.
- ▶ Mills (1979): If D is an \aleph_1 -algebraic distributive lattice, and $\mathcal{M} \models \text{PA}$, there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong D$.

- ▶ Gaifman (1960s, published 1976): Every model $\mathcal{M} \models \text{PA}$ has a **minimal** elementary end extension
- ▶ Blass (1974): Every countable nonstandard \mathcal{M} has a minimal cofinal extension.

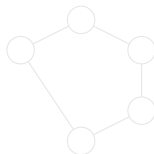
Definition

A lattice L is called **distributive** if whenever $a, b, c \in L$,
 $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$.

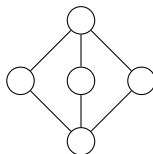
- ▶ Gaifman: Given a finite **distributive** lattice D , there is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong D$.
- ▶ Mills (1979): If D is an \aleph_1 -algebraic distributive lattice, and $\mathcal{M} \models \text{PA}$, there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong D$.



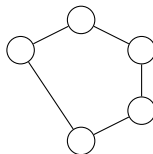
- ▶ Paris (1977): There is $\mathcal{M} \models \text{PA}$ such that the lattice \mathbf{M}_3 is isomorphic to a sublattice of $\text{Lt}(\mathcal{M})$.
- ▶ Schmerl (1986): There is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong \mathbf{M}_3$.



- ▶ Wilkie (1977): There is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong \mathbf{N}_5$. In particular, we can take an \mathcal{M} such that $\mathbb{N} \prec \mathcal{M}$.



- ▶ Paris (1977): There is $\mathcal{M} \models \text{PA}$ such that the lattice \mathbf{M}_3 is isomorphic to a sublattice of $\text{Lt}(\mathcal{M})$.
- ▶ Schmerl (1986): There is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong \mathbf{M}_3$.



- ▶ Wilkie (1977): There is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong \mathbf{N}_5$. In particular, we can take an \mathcal{M} such that $\mathbb{N} \prec \mathcal{M}$.

The Lattice Problem

Question

For which lattices L is it the case that there is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong L$?

Question

Given a model $\mathcal{M} \models \text{PA}$, for which lattices L is it the case that there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$?

Theorems are often of the form: “Every (countable)/(nonstandard) \mathcal{M} has an elementary (end/cofinal) extension \mathcal{N} such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$ ”.

Question

For which lattices L is it the case that there is $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) \cong L$?

Question

Given a model $\mathcal{M} \models \text{PA}$, for which lattices L is it the case that there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$?

Theorems are often of the form: “Every (countable)/(nonstandard) \mathcal{M} has an elementary (end/cofinal) extension \mathcal{N} such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$ ”.

The Lattice Problem, II

We also are interested in knowing the nature of these extensions: are they cofinal extensions, end extensions, can they be conservative, etc.

Example

There is no $\mathcal{M} \models \text{TA}$ such that $\text{Lt}(\mathcal{M}) \cong \mathbb{M}_3$.

Example

It is unknown if there are any models $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) = \mathbb{M}_{16}$.

The Lattice Problem, II

We also are interested in knowing the nature of these extensions: are they cofinal extensions, end extensions, can they be conservative, etc.

Example

There is no $\mathcal{M} \models \text{TA}$ such that $\text{Lt}(\mathcal{M}) \cong \mathbf{M}_3$.

Example

It is unknown if there are any models $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) = \mathbf{M}_{16}$.

The Lattice Problem, II

We also are interested in knowing the nature of these extensions: are they cofinal extensions, end extensions, can they be conservative, etc.

Example

There is no $\mathcal{M} \models \text{TA}$ such that $\text{Lt}(\mathcal{M}) \cong \mathbf{M}_3$.

Example

It is unknown if there are any models $\mathcal{M} \models \text{PA}$ such that $\text{Lt}(\mathcal{M}) = \mathbf{M}_{16}$.

Representations of Lattices

Many proofs about interstructure lattices involve representations of lattices.

Definition

Let A be any set. Then $\text{Eq}(A)$ is the set of equivalence relations on that set, and it forms a lattice under inclusion. $\mathbf{0}_A$ is the discrete relation, and $\mathbf{1}_A$ is the trivial relation.

Definition

Let L be a finite lattice. $\alpha : L \rightarrow \text{Eq}(A)$ is called a **representation** of L on A if:

1. α is an injection,
2. $\alpha(0) = \mathbf{1}_A$ and $\alpha(1) = \mathbf{0}_A$, and
3. For all $a, b \in L$, $\alpha(a \vee b) = \alpha(a) \wedge \alpha(b)$.

If, in addition, we have, for all $a, b \in L$, $\alpha(a \wedge b) = \alpha(a) \vee \alpha(b)$, then we say α is a **lattice representation**.

Representations of Lattices

Many proofs about interstructure lattices involve representations of lattices.

Definition

Let A be any set. Then $\text{Eq}(A)$ is the set of equivalence relations on that set, and it forms a lattice under inclusion. $\mathbf{0}_A$ is the discrete relation, and $\mathbf{1}_A$ is the trivial relation.

Definition

Let L be a finite lattice. $\alpha : L \rightarrow \text{Eq}(A)$ is called a **representation** of L on A if:

1. α is an injection,
2. $\alpha(0) = \mathbf{1}_A$ and $\alpha(1) = \mathbf{0}_A$, and
3. For all $a, b \in L$, $\alpha(a \vee b) = \alpha(a) \wedge \alpha(b)$.

If, in addition, we have, for all $a, b \in L$, $\alpha(a \wedge b) = \alpha(a) \vee \alpha(b)$, then we say α is a **lattice representation**.

Representations of Lattices

Many proofs about interstructure lattices involve representations of lattices.

Definition

Let A be any set. Then $\text{Eq}(A)$ is the set of equivalence relations on that set, and it forms a lattice under inclusion. $\mathbf{0}_A$ is the discrete relation, and $\mathbf{1}_A$ is the trivial relation.

Definition

Let L be a finite lattice. $\alpha : L \rightarrow \text{Eq}(A)$ is called a **representation** of L on A if:

1. α is an injection,
2. $\alpha(0) = \mathbf{1}_A$ and $\alpha(1) = \mathbf{0}_A$, and
3. For all $a, b \in L$, $\alpha(a \vee b) = \alpha(a) \wedge \alpha(b)$.

If, in addition, we have, for all $a, b \in L$, $\alpha(a \wedge b) = \alpha(a) \vee \alpha(b)$, then we say α is a **lattice representation**.

Representations of Lattices

Many proofs about interstructure lattices involve representations of lattices.

Definition

Let A be any set. Then $\text{Eq}(A)$ is the set of equivalence relations on that set, and it forms a lattice under inclusion. $\mathbf{0}_A$ is the discrete relation, and $\mathbf{1}_A$ is the trivial relation.

Definition

Let L be a finite lattice. $\alpha : L \rightarrow \text{Eq}(A)$ is called a **representation** of L on A if:

1. α is an injection,
2. $\alpha(0) = \mathbf{1}_A$ and $\alpha(1) = \mathbf{0}_A$, and
3. For all $a, b \in L$, $\alpha(a \vee b) = \alpha(a) \wedge \alpha(b)$.

If, in addition, we have, for all $a, b \in L$, $\alpha(a \wedge b) = \alpha(a) \vee \alpha(b)$, then we say α is a **lattice representation**.

Example: 2

Let $\mathcal{M} \models \text{PA}$ be countable, $A \subseteq M$ be unbounded and definable, and $\alpha : \mathbf{2} \rightarrow \text{Eq}(A)$ be the (unique) representation.

If $\Theta \in \text{Eq}(M)$ is definable, then there is an unbounded, definable $B \subseteq A$ such that either $\Theta \cap B^2 = \alpha(0) \cap B^2$ or $\Theta \cap B^2 = \alpha(1) \cap B^2$. (Every definable Skolem function can be made 1 to 1 or constant on an unbounded set.)

Enumerate the definable equivalence relations, apply the above statement to obtain a sequence $M = A_0 \supseteq A_1 \supseteq A_2 \supseteq \dots$

Let $p(x) = \{\phi(x) \mid \exists n < \omega \mathcal{M} \models \forall x \in A_n \phi(x)\}$. If c realizes $p(x)$, then $\mathcal{M} \prec_{\text{end}} \mathcal{M}(c)$ and $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong \mathbf{2}$.

Example: 2

Let $\mathcal{M} \models \text{PA}$ be countable, $A \subseteq M$ be unbounded and definable, and $\alpha : \mathbf{2} \rightarrow \text{Eq}(A)$ be the (unique) representation.

If $\Theta \in \text{Eq}(M)$ is definable, then there is an unbounded, definable $B \subseteq A$ such that either $\Theta \cap B^2 = \alpha(0) \cap B^2$ or $\Theta \cap B^2 = \alpha(1) \cap B^2$. (Every definable Skolem function can be made 1 to 1 or constant on an unbounded set.)

Enumerate the definable equivalence relations, apply the above statement to obtain a sequence $M = A_0 \supseteq A_1 \supseteq A_2 \supseteq \dots$

Let $p(x) = \{\phi(x) \mid \exists n < \omega \mathcal{M} \models \forall x \in A_n \phi(x)\}$. If c realizes $p(x)$, then $\mathcal{M} \prec_{\text{end}} \mathcal{M}(c)$ and $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong \mathbf{2}$.

Example: 2

Let $\mathcal{M} \models \text{PA}$ be countable, $A \subseteq M$ be unbounded and definable, and $\alpha : \mathbf{2} \rightarrow \text{Eq}(A)$ be the (unique) representation.

If $\Theta \in \text{Eq}(M)$ is definable, then there is an unbounded, definable $B \subseteq A$ such that either $\Theta \cap B^2 = \alpha(0) \cap B^2$ or $\Theta \cap B^2 = \alpha(1) \cap B^2$. (Every definable Skolem function can be made 1 to 1 or constant on an unbounded set.)

Enumerate the definable equivalence relations, apply the above statement to obtain a sequence $M = A_0 \supseteq A_1 \supseteq A_2 \supseteq \dots$

Let $p(x) = \{\phi(x) \mid \exists n < \omega \mathcal{M} \models \forall x \in A_n \phi(x)\}$. If c realizes $p(x)$, then $\mathcal{M} \prec_{\text{end}} \mathcal{M}(c)$ and $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong \mathbf{2}$.

Example: 2

Let $\mathcal{M} \models \text{PA}$ be countable, $A \subseteq M$ be unbounded and definable, and $\alpha : \mathbf{2} \rightarrow \text{Eq}(A)$ be the (unique) representation.

If $\Theta \in \text{Eq}(M)$ is definable, then there is an unbounded, definable $B \subseteq A$ such that either $\Theta \cap B^2 = \alpha(0) \cap B^2$ or $\Theta \cap B^2 = \alpha(1) \cap B^2$. (Every definable Skolem function can be made 1 to 1 or constant on an unbounded set.)

Enumerate the definable equivalence relations, apply the above statement to obtain a sequence $M = A_0 \supseteq A_1 \supseteq A_2 \supseteq \dots$

Let $p(x) = \{\phi(x) \mid \exists n < \omega \mathcal{M} \models \forall x \in A_n \phi(x)\}$. If c realizes $p(x)$, then $\mathcal{M} \prec_{\text{end}} \mathcal{M}(c)$ and $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong \mathbf{2}$.

Representations of Lattices

Definition

Let $\alpha : L \rightarrow \text{Eq}(A)$ be a representation, and suppose $B \subseteq A$. Then $\alpha|B : L \rightarrow \text{Eq}(B)$ is the function defined by $(\alpha|B)(r) = \alpha(r) \cap B^2$ for each $r \in L$.

Remark

$\alpha|B$ is not always a representation (for example, let $A = \mathbb{N}$, and $B = \{0\}$, then $(\alpha|B)(1) = (\alpha|B)(0)$).

Definition

Let L be a finite lattice, and $\alpha : L \rightarrow \text{Eq}(A)$ a representation. We say α is **0-CPP** if for each $r > 0$, $\alpha(r)$ has more than 2 classes. We say α is **$(n+1)$ -CPP** if for each $\Theta \in \text{Eq}(A)$, there is $B \subseteq A$ such that $\alpha|B$ is an n -CPP representation and $r \in L$ such that $\Theta \cap B^2 = \alpha(r) \cap B^2$.

Definition

Let $\alpha : L \rightarrow \text{Eq}(A)$ be a representation, and suppose $B \subseteq A$. Then $\alpha|B : L \rightarrow \text{Eq}(B)$ is the function defined by $(\alpha|B)(r) = \alpha(r) \cap B^2$ for each $r \in L$.

Remark

$\alpha|B$ is not always a representation (for example, let $A = \mathbb{N}$, and $B = \{0\}$, then $(\alpha|B)(1) = (\alpha|B)(0)$).

Definition

Let L be a finite lattice, and $\alpha : L \rightarrow \text{Eq}(A)$ a representation. We say α is **0-CPP** if for each $r > 0$, $\alpha(r)$ has more than 2 classes. We say α is **$(n+1)$ -CPP** if for each $\Theta \in \text{Eq}(A)$, there is $B \subseteq A$ such that $\alpha|B$ is an n -CPP representation and $r \in L$ such that $\Theta \cap B^2 = \alpha(r) \cap B^2$.

Definition

Let $\alpha : L \rightarrow \text{Eq}(A)$ be a representation, and suppose $B \subseteq A$. Then $\alpha|B : L \rightarrow \text{Eq}(B)$ is the function defined by $(\alpha|B)(r) = \alpha(r) \cap B^2$ for each $r \in L$.

Remark

$\alpha|B$ is not always a representation (for example, let $A = \mathbb{N}$, and $B = \{0\}$, then $(\alpha|B)(1) = (\alpha|B)(0)$).

Definition

Let L be a finite lattice, and $\alpha : L \rightarrow \text{Eq}(A)$ a representation. We say α is **0-CPP** if for each $r > 0$, $\alpha(r)$ has more than 2 classes. We say α is **$(n+1)$ -CPP** if for each $\Theta \in \text{Eq}(A)$, there is $B \subseteq A$ such that $\alpha|B$ is an n -CPP representation and $r \in L$ such that $\Theta \cap B^2 = \alpha(r) \cap B^2$.

Remark

If L has an n -CPP representation, then it has a finite n -CPP representation. This can be expressed in PA with a Σ_1 formula $cpp(L, n)$.

Theorem

Let \mathcal{M} be a countable nonstandard model of PA^ and L a finite lattice. If $\mathcal{M} \models cpp(L, n)$ for each $n \in \omega$, then \mathcal{M} has a **cofinal** extension \mathcal{N} such that $Lt(\mathcal{N}/\mathcal{M}) \cong L$.*

Theorem

Let L be a finite lattice, $\mathcal{M} \models PA^$, and $Lt(\mathcal{N}/\mathcal{M}) \cong L$. Then $\mathcal{M} \models cpp(L, n)$ for each $n \in \omega$.*

Remark

If L has an n -CPP representation, then it has a finite n -CPP representation. This can be expressed in PA with a Σ_1 formula $cpp(L, n)$.

Theorem

Let \mathcal{M} be a countable nonstandard model of PA^* and L a finite lattice. If $\mathcal{M} \models cpp(L, n)$ for each $n \in \omega$, then \mathcal{M} has a *cofinal* extension \mathcal{N} such that $Lt(\mathcal{N}/\mathcal{M}) \cong L$.

Theorem

Let L be a finite lattice, $\mathcal{M} \models PA^*$, and $Lt(\mathcal{N}/\mathcal{M}) \cong L$. Then $\mathcal{M} \models cpp(L, n)$ for each $n \in \omega$.

Remark

If L has an n -CPP representation, then it has a finite n -CPP representation. This can be expressed in PA with a Σ_1 formula $cpp(L, n)$.

Theorem

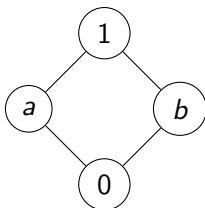
Let \mathcal{M} be a countable nonstandard model of PA^* and L a finite lattice. If $\mathcal{M} \models cpp(L, n)$ for each $n \in \omega$, then \mathcal{M} has a *cofinal* extension \mathcal{N} such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$.

Theorem

Let L be a finite lattice, $\mathcal{M} \models \text{PA}^*$, and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$. Then $\mathcal{M} \models cpp(L, n)$ for each $n \in \omega$.

Theorem

Let $\mathcal{M} \models \text{PA}$ be countable (standard or non-standard). Then \mathcal{M} has an elementary end extension \mathcal{N} such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong B_2$.



Representation of \mathbf{B}_2

Let $A = \{(x, y) \mid x \geq y\}$. Define a representation $\alpha : \mathbf{B}_2 \rightarrow \text{Eq}(A)$ by letting $\langle (x_0, y_0), (x_1, y_1) \rangle \in \alpha(a)$ iff $x_0 = x_1$, and $\langle (x_0, y_0), (x_1, y_1) \rangle \in \alpha(b)$ iff $y_0 = y_1$.

We need a combinatorial lemma, similar to the one for the construction of a minimal extension.

Lemma

Let $\Theta \in \text{Eq}(A)$ be definable. Then there is $B \subseteq A$ such that $\alpha|_B \cong \alpha$ and $r \in \mathbf{B}_2$ such that $\Theta \cap B^2 = \alpha(r) \cap B^2$.

Representation of \mathbf{B}_2

Let $A = \{(x, y) \mid x \geq y\}$. Define a representation $\alpha : \mathbf{B}_2 \rightarrow \text{Eq}(A)$ by letting $\langle (x_0, y_0), (x_1, y_1) \rangle \in \alpha(a)$ iff $x_0 = x_1$, and $\langle (x_0, y_0), (x_1, y_1) \rangle \in \alpha(b)$ iff $y_0 = y_1$.

We need a combinatorial lemma, similar to the one for the construction of a minimal extension.

Lemma

Let $\Theta \in \text{Eq}(A)$ be definable. Then there is $B \subseteq A$ such that $\alpha|_B \cong \alpha$ and $r \in \mathbf{B}_2$ such that $\Theta \cap B^2 = \alpha(r) \cap B^2$.

Representation of \mathbf{B}_2

Let $A = \{(x, y) \mid x \geq y\}$. Define a representation $\alpha : \mathbf{B}_2 \rightarrow \text{Eq}(A)$ by letting $\langle (x_0, y_0), (x_1, y_1) \rangle \in \alpha(a)$ iff $x_0 = x_1$, and $\langle (x_0, y_0), (x_1, y_1) \rangle \in \alpha(b)$ iff $y_0 = y_1$.

We need a combinatorial lemma, similar to the one for the construction of a minimal extension.

Lemma

Let $\Theta \in \text{Eq}(A)$ be definable. Then there is $B \subseteq A$ such that $\alpha|_B \cong \alpha$ and $r \in \mathbf{B}_2$ such that $\Theta \cap B^2 = \alpha(r) \cap B^2$.

Representation of \mathbb{B}_2

Enumerate all the definable equivalence relations. Starting with $A_0 = A$, we obtain a sequence of definable sets $A_0 \supseteq A_1 \supseteq \dots$ and we can construct a type

$$p(x) = \{\phi(x) \mid \exists n < \omega \mathcal{M} \models \forall x \in A_n \phi(x)\}$$

For $r \in \mathbb{B}_2$, let t_r be the Skolem term mapping x to the least y such that $\langle x, y \rangle \in \alpha(r)$.

Let c realize $p(x)$ and $\mathcal{N} = \mathcal{M}(c)$. Then $\mathcal{M} \prec_{\text{end}} \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong \mathbb{B}_2$, with the other submodels \mathcal{M}_a and \mathcal{M}_b being generated by $t_a(c)$ and $t_b(c)$, respectively.

Representation of \mathbb{B}_2

Enumerate all the definable equivalence relations. Starting with $A_0 = A$, we obtain a sequence of definable sets $A_0 \supseteq A_1 \supseteq \dots$ and we can construct a type

$$p(x) = \{\phi(x) \mid \exists n < \omega \mathcal{M} \models \forall x \in A_n \phi(x)\}$$

For $r \in \mathbb{B}_2$, let t_r be the Skolem term mapping x to the least y such that $\langle x, y \rangle \in \alpha(r)$.

Let c realize $p(x)$ and $\mathcal{N} = \mathcal{M}(c)$. Then $\mathcal{M} \prec_{\text{end}} \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong \mathbb{B}_2$, with the other submodels \mathcal{M}_a and \mathcal{M}_b being generated by $t_a(c)$ and $t_b(c)$, respectively.

Representation of \mathbf{B}_2

Enumerate all the definable equivalence relations. Starting with $A_0 = A$, we obtain a sequence of definable sets $A_0 \supseteq A_1 \supseteq \dots$ and we can construct a type

$$p(x) = \{\phi(x) \mid \exists n < \omega \mathcal{M} \models \forall x \in A_n \phi(x)\}$$

For $r \in \mathbf{B}_2$, let t_r be the Skolem term mapping x to the least y such that $\langle x, y \rangle \in \alpha(r)$.

Let c realize $p(x)$ and $\mathcal{N} = \mathcal{M}(c)$. Then $\mathcal{M} \prec_{\text{end}} \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong \mathbf{B}_2$, with the other submodels \mathcal{M}_a and \mathcal{M}_b being generated by $t_a(c)$ and $t_b(c)$, respectively.

Theorem (Schmerl 2014)

If \mathcal{M}, \mathcal{N} are countable, arithmetically saturated models of PA, and $\text{Aut}(\mathcal{M}) \cong \text{Aut}(\mathcal{N})$, then $\text{Th}(\mathcal{M})' \equiv_T \text{Th}(\mathcal{N})'$.

The proof is very involved, so we mention part of one step here. To $X \subseteq \omega$, we associate an infinite distributive lattice $\mathcal{D}'(X)$.

Lemma

Let \mathcal{M} be recursively saturated. Then there is $\mathcal{M}_1 \prec \mathcal{M}$ such that $\text{Lt}(\mathcal{M}_1) \cong \mathcal{D}'(X)$ iff there is $Y \in \text{SSy}(\mathcal{M})$ such that $X \leq_T Y'$.

Using this (and a lot of power from the automorphism groups), we can then show that $\text{Th}(\mathcal{M})' \leq_T \text{Th}(\mathcal{N})'$.

Theorem (Schmerl 2014)

If \mathcal{M}, \mathcal{N} are countable, arithmetically saturated models of PA, and $\text{Aut}(\mathcal{M}) \cong \text{Aut}(\mathcal{N})$, then $\text{Th}(\mathcal{M})' \equiv_T \text{Th}(\mathcal{N})'$.

The proof is very involved, so we mention part of one step here. To $X \subseteq \omega$, we associate an infinite distributive lattice $\mathcal{D}'(X)$.

Lemma

Let \mathcal{M} be recursively saturated. Then there is $\mathcal{M}_1 \prec \mathcal{M}$ such that $\text{Lt}(\mathcal{M}_1) \cong \mathcal{D}'(X)$ iff there is $Y \in \text{SSy}(\mathcal{M})$ such that $X \leq_T Y'$.

Using this (and a lot of power from the automorphism groups), we can then show that $\text{Th}(\mathcal{M})' \leq_T \text{Th}(\mathcal{N})'$.

Theorem (Schmerl 2014)

If \mathcal{M}, \mathcal{N} are countable, arithmetically saturated models of PA, and $\text{Aut}(\mathcal{M}) \cong \text{Aut}(\mathcal{N})$, then $\text{Th}(\mathcal{M})' \equiv_T \text{Th}(\mathcal{N})'$.

The proof is very involved, so we mention part of one step here. To $X \subseteq \omega$, we associate an infinite distributive lattice $\mathcal{D}'(X)$.

Lemma

Let \mathcal{M} be recursively saturated. Then there is $\mathcal{M}_1 \prec \mathcal{M}$ such that $\text{Lt}(\mathcal{M}_1) \cong \mathcal{D}'(X)$ iff there is $Y \in \text{SSy}(\mathcal{M})$ such that $X \leq_T Y'$.

Using this (and a lot of power from the automorphism groups), we can then show that $\text{Th}(\mathcal{M})' \leq_T \text{Th}(\mathcal{N})'$.

Theorem (Schmerl 2014)

If \mathcal{M}, \mathcal{N} are countable, arithmetically saturated models of PA, and $\text{Aut}(\mathcal{M}) \cong \text{Aut}(\mathcal{N})$, then $\text{Th}(\mathcal{M})' \equiv_T \text{Th}(\mathcal{N})'$.

The proof is very involved, so we mention part of one step here. To $X \subseteq \omega$, we associate an infinite distributive lattice $\mathcal{D}'(X)$.

Lemma

Let \mathcal{M} be recursively saturated. Then there is $\mathcal{M}_1 \prec \mathcal{M}$ such that $\text{Lt}(\mathcal{M}_1) \cong \mathcal{D}'(X)$ iff there is $Y \in \text{SSy}(\mathcal{M})$ such that $X \leq_T Y'$.

Using this (and a lot of power from the automorphism groups), we can then show that $\text{Th}(\mathcal{M})' \leq_T \text{Th}(\mathcal{N})'$.