

# Ehrenfeucht principles in set theory

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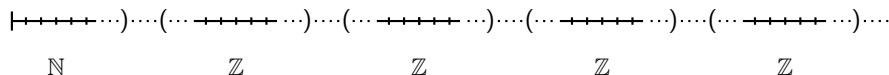
## Models of arithmetic

### Peano Arithmetic

- The first-order language of arithmetic is  $\mathcal{L}_A = \{+, \cdot, <, 0, 1\}$ .
- **Peano Arithmetic (PA)** is a collection of statements in  $\mathcal{L}_A$  codifying the **fundamental properties of the natural numbers**.
  - ▶ Commutativity and associativity of addition and multiplication, distributive law, ordering is discrete with least element 0, 0 is the additive identity, etc.
  - ▶ **Induction scheme**: for every  $\mathcal{L}_A$ -formula  $\varphi(x, \vec{y})$ ,

$$\forall \vec{y} [(\varphi(0, \vec{y}) \wedge (\forall x \varphi(x, \vec{y}) \rightarrow \varphi(x+1, \vec{y}))) \rightarrow \forall x \varphi(x, \vec{y})].$$

- PA proves the **least-number principle**: every definable (with parameters) set has a  $<$ -least element.
- The **standard** model of PA is  $\mathbb{N}$ .
- A **countable nonstandard** model of PA looks like:

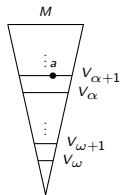


densely many copies of  $\mathbb{Z}$

## Models of set theory

### Zermelo-Fraenkel set theory (with choice)

- The first-order language of set theory is  $\mathcal{L}_S = \{\in\}$ .
- **Zermelo-Fraenkel set theory (ZF(C))** is a collection of statements in  $\mathcal{L}_S$  codifying the **fundamental properties of sets**.
- A model  $M \models \text{ZF}$  is the union of the **von Neumann hierarchy**  $M = \bigcup_{\alpha \in \text{ORD}} V_\alpha^M$ :
  - ▶  $V_0^M = \emptyset$ ,
  - ▶  $V_{\alpha+1}^M = \mathcal{P}^M(V_\alpha)$  (Powerset of  $V_\alpha^M$ ),
  - ▶  $V_\lambda^M = \bigcup_{\alpha < \lambda} V_\alpha^M$  for limit ordinals  $\lambda$ .
  - ▶ If  $a \in M$ , then **rank**( $a$ ) =  $\alpha$  is **least** such that  $a \in V_{\alpha+1}^M$ .
- **Theorem:** (Lévy-Montague reflection) For every  $n$ , there are unboundedly many ordinals  $\alpha$  in  $M$  such that  $V_\alpha^M \prec_{\Sigma_n} M$ .



There is an active **exchange of concepts, methods and techniques** between model theory of models of PA and model theory of models of ZF.

## Some general model theory

Suppose  $M$  is a first-order structure and  $P \subseteq M$ .

### Definition:

- $a \neq b$  in  $M$  are **indiscernible** if they have the **same type**:  
for every  $\varphi(x)$ ,  $M \models \varphi(a) \leftrightarrow \varphi(b)$ .
- $a \neq b$  in  $M$  are  **$P$ -indiscernible** if they have the same type with **parameters from  $P$** .
- Otherwise they are  **$(P)$ -discernible**.

It is **common** for models to have indiscernible elements:

- A first-order structure (in a countable language) of **size greater than  $2^\omega$**  has indiscernible elements.
- A **non-rigid** first-order structure has indiscernible elements.
- If  $T$  is a first-order theory with an infinite model, then there are models of  $T$  with sets of indiscernible elements of any cardinality (by compactness).

## Ehrenfeucht's lemma in arithmetic

A powerful tool in model theory of models of PA is

**Ehrenfeucht's lemma:** (Ehrenfeucht, '73) If  $a \neq b$  in a model  $M \models \text{PA}$  and  $b$  is definable from  $a$  in  $M$ , then  $a$  and  $b$  are discernible in  $M$ .

Here is an application.

**Theorem:** If  $M \models \text{PA}$  is the Skolem closure of a single element  $a \in M$ , then it has no non-trivial automorphisms.

**Proof:**

- $a$  is the only element of its type by Ehrenfeucht's lemma.
- An automorphism of  $M$  must take  $a$  to  $a$  and therefore fixes everything.  $\square$

**Note:** Given any first-order structure, we can ask if Ehrenfeucht's lemma holds there.

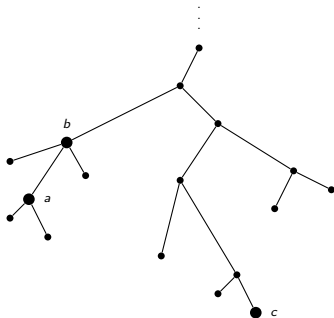
**Question:** Does Ehrenfeucht's lemma hold for models of ZF?

# Proof of Ehrenfeucht's lemma

**Proof:** Suppose  $a \neq b$  and  $b$  is definable from  $a$ .  
Fix a definable function  $f(x)$  such that  $f(a) = b$ .

**Case 1:**  $a < f(a) = b$

- $G$  is the graph whose edges are pairs  $\{x, f(x)\}$  such that  $x < f(x)$ .
- $G$  has an edge between  $a$  and  $b$ .
- $G$  is loop-free.
- $d(x, y)$  is the length of the shortest path between  $x$  and  $y$  in  $G$ , if connected.
- $c$  is  $<$ -least in the connected component of  $a$  and  $b$  (exists by the least-number principle).
- $d(c, a)$  and  $d(c, b)$  differ by 1.
- $d(c, a)$  is even iff  $d(c, b)$  is odd.

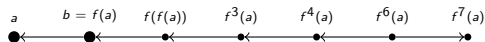


## Proof of Ehrenfeucht's lemma (continued)

$f(x)$  is a definable function such that  $f(a) = b$ .

**Case 2:**  $a > f(a) = b$

- $d(x)$  is the number of times  $f(x)$  can be iterated before the values start to **increase**.
- $d(x)$  exists by the least-number principle.
- $d(a)$  and  $d(b)$  differ by 1.



□



## Ordinal definable sets

Suppose  $M \models \text{ZF}$  and fix a definable bijection  $F : \text{ORD}^{<\omega^M} \xrightarrow[\text{onto}]{1-1} \text{ORD}$ .

**Definition:**

- A set is **ordinal definable** if it is definable with ordinal parameters (using  $F$ , we can always assume that there is a **single** ordinal parameter).
- **OD** is the collection of all ordinal definable sets.
- **HOD** is the collection of all hereditarily ordinal definable sets.

**Lemma:** A set  $a \in \text{OD}$  iff  $M$  satisfies that  $a$  is ordinal definable in some  $V_\alpha$ .

**Proof:**

( $\Rightarrow$ ): If  $a \in \text{OD}$ , then  $a$  is defined by the same formula in some  $V_\alpha$  (by reflection).

( $\Leftarrow$ ): Suppose  $M$  satisfies that  $a$  is defined by  $\ulcorner \varphi(x, \beta) \urcorner$  in  $V_\alpha$ .

Note that  $\varphi$  might be **nonstandard**. But

$$\psi(x, \ulcorner \varphi \urcorner, \beta, \alpha) := \exists y y = V_\alpha \wedge y \models \ulcorner \varphi(x, \beta) \urcorner$$

defines  $a$  in  $M$ .  $\square$

**Corollary:** The collection **OD** is **first-order definable** in  $M$ .

## A definable well-ordering of OD

**Lemma:**  $M$  has a definable well-ordering  $<$  of OD.

**Proof:** Fix  $x, y \in \text{OD}$ .

- $x$  is definable in  $V_\alpha$  by  $\ulcorner \varphi(x, \xi) \urcorner$ , where  $\alpha$  and  $F(\ulcorner \varphi \urcorner, \xi)$  are least.
  - $y$  is definable in  $V_\beta$  by  $\ulcorner \psi(x, \mu) \urcorner$ , where  $\beta$  and  $F(\ulcorner \psi \urcorner, \mu)$  are least.
- $x < y$  if  $\alpha < \beta$  or  $\alpha = \beta$  and  $F(\ulcorner \varphi \urcorner, \xi) < F(\ulcorner \psi \urcorner, \mu)$ .  $\square$

**Lemma:** HOD is a transitive model of ZFC.

**Proof:**

- Transitivity is by definition.
- ZF a is direct verification.
- If  $a \in \text{HOD}$ , then the well-ordering of OD restricted to  $a$  is ordinal definable.

## The axiom: $V = \text{HOD}$

**Definition:** The axiom  $V = \text{HOD}$  states that every set is hereditarily ordinal definable.

**Corollary:** If  $M \models \text{ZF}$ , then  $M$  is a model of  $V = \text{HOD}$  iff  $M$  has a **definable well-ordering**.

**Proof:**

( $\Rightarrow$ ):  $M$  has a definable well-ordering of ordinal definable sets.

( $\Leftarrow$ ): Assume (wlog) order-type of well-ordering is ORD. Each  $a \in M$  is the  $\alpha^{\text{th}}$ -element for some ordinal  $\alpha$ .  $\square$

## Ehrenfeucht's lemma and ordinal definable sets

**Theorem:** If  $a \neq b$  in a model  $M \models \text{ZF}$ ,  $a$  is ordinal definable and  $b$  is definable from  $a$  in  $M$ , then  $a$  and  $b$  are discernible in  $M$ .

“Ehrenfeucht's lemma holds for ordinal definable sets.”

**Proof:**

- $b$  is ordinal definable.
- Use same argument as for models of PA with the well-ordering  $<$  of OD-sets.  $\square$

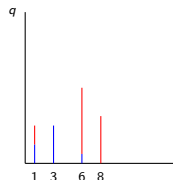
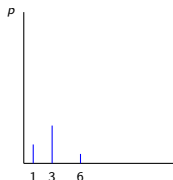
**Corollary:** Ehrenfeucht's lemma holds for every model of  $\text{ZF} + V = \text{HOD}$ .

## A counterexample to Ehrenfeucht's lemma

**Theorem:** (Fuchs, G., Hamkins) If  $M \models \text{ZFC}$  and  $M[c]$  is the Cohen forcing extension, then there are  $a \neq b$  in  $M[c]$  such that  $a$  and  $b$  are inter-definable, but  $a$  and  $b$  are  $M$ -indiscernible.

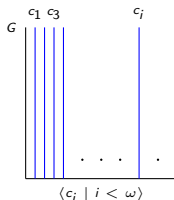
**Proof:**

- Poset  $\mathbb{P}$  to add a Cohen real is isomorphic to finite-support product  $\prod_{i < \omega} \mathbb{P}$ .
- The poset  $\prod_{i < \omega} \mathbb{P}$  adds  $\omega$ -many Cohen reals:
  - ▶ conditions: finite functions  $p : \text{dom}(p) \rightarrow {}^{<\omega}2$  on  $\omega$ ,
  - ▶ order:  $q \leq p$  if  $\text{dom}(p) \subseteq \text{dom}(q)$  and for all  $n \in \text{dom}(p)$ ,  $q(n)$  extends  $p(n)$ .



## A counterexample to Ehrenfeucht's lemma (continued)

- $G \subseteq \prod_{i < \omega} \mathbb{P}$  is  $M$ -generic.
- $C = \{c_i \mid i < \omega\}$ .
- $E = \{c_{2i} \mid i < \omega\}$  and  $O = \{c_{2i+1} \mid i < \omega\}$ .
- $\langle C, E \rangle$  and  $\langle C, O \rangle$  are  $M$ -indiscernible.



- $\dot{C}$ ,  $\dot{E}$ , and  $\dot{O}$  are  $\mathbb{P}$ -names for  $C$ ,  $E$ , and  $O$ .
- Suppose  $M[G] \models \varphi(\langle C, E \rangle, d)$  with  $d \in M$ . Then there is  $q \Vdash \varphi(\langle \dot{C}, \dot{E} \rangle, \check{d})$  in  $G$ .
- If  $i \in \text{dom}(q)$  is even/odd, then there is odd/even  $m_i$  such that  $q(i) \subseteq c_{m_i}$ .
- $\pi : \mathbb{P} \rightarrow \mathbb{P}$  is an automorphism such that
  - ▶  $\pi$  switches even and odd coordinates.
  - ▶  $\pi(i) = m_i$  for all  $i \in \text{dom}(q)$ .
- $H = \pi " G$  is  $M$ -generic for  $\prod_{i < \omega} \mathbb{P}$  and  $M[G] = M[H]$ .
- $q \in H$ .
- $\dot{C}_H = C$ ,  $\dot{E}_H = O$ , and  $\dot{O}_H = E$ .
- $M[H] = M[G] \models \varphi(\langle C, O \rangle, d)$ .  $\square$

# Is Ehrenfeucht's lemma equivalent to $V = \text{HOD}$ ?

**Question:** If Ehrenfeucht's lemma holds in  $M \models \text{ZF}(C)$ , is  $M$  a model of  $V = \text{HOD}$ ?

# Ehrenfeucht Principles

Parametric generalizations of Ehrenfeucht's lemma.

**Definition:** (Fuchs, G., Hamkins) Let  $M \models \text{ZF}$  and let  $A, P, Q \subseteq M$ . The principle  $\text{EL}(A, P, Q)$  for  $M$  asserts that if  $a \in A$ ,  $a \neq b$ , and  $b$  is definable in  $M$  from  $a$  with parameters from  $P$ , then  $a$  and  $b$  are  $Q$ -discernible.

" $P$ -definability from  $A$  implies  $Q$ -discernibility."

**Observation:**

- $\text{EL}(M, \emptyset, \emptyset)$  is Ehrenfeucht's lemma.
- $\text{EL}(\text{OD}, \emptyset, \emptyset)$  holds for every  $M \models \text{ZF}$ .
- $\text{EL}(M[c], \emptyset, M)$  fails in every Cohen forcing extension  $M[c]$ .
- $\text{EL}(M, \emptyset, \emptyset)$  implies  $\text{EL}(M, P, P)$  for every  $P$ .
- $\text{EL}(A, P, Q)$  gets stronger if  $A, P$  are enlarged or  $Q$  is diminished.



## Interesting Ehrenfeucht principles

Suppose  $M \models \text{ZF}$ .

$\text{EL}(M, \text{ORD}, \text{ORD})$

- If  $a \neq b$  and  $b$  is definable from  $a$  with **ordinal parameters**, then  $a$  and  $b$  are **ORD-discernible**.
- **First-order** expressible (similar argument to definability of OD).

$\text{EL}(M, M, \emptyset)$

- $M$  contains **no indiscernibles**.
- (Leibniz) **Identity of Indiscernibles**: any two objects differ on some property.
- (Enayat)  $M$  is **Leibnizian**.
- **Theorem**: (Enayat, '04) There are **uncountable** Leibnizian models of set theory.
- **Theorem**: (Mycielski, '95) A theory  $T$  extending ZF has a **Leibnizian model** iff  $T$  **proves the Leibniz-Mycielski axiom** (next slide).

$\text{EL}(M, M, \text{ORD})$

- Any  $a \neq b$  are **ORD-discernible**.
- **First-order** expressible.

## Leibniz-Mycielski axiom

**Leibniz-Mycielski axiom (LM):** (Mycielski, '95) If  $a \neq b$  in  $M \models \text{ZF}$ , then  $a$  and  $b$  are discernible in some  $V_\alpha$ .

**Observation:** Suppose  $M \models \text{ZF}$ .

- LM is equivalent to  $\text{EL}(M, M, \text{ORD})$ .

**Proof:** Fix  $a \neq b$ .

$(\Rightarrow)$ : Suppose LM holds.

- ▶  $V_\alpha \models \varphi(a) \wedge \neg\varphi(b)$ .
- ▶  $\psi(x, \alpha) := \exists y y = V_\alpha \wedge y \models \varphi(x)$ .
- ▶  $M \models \psi(a, \alpha) \wedge \neg\psi(b, \alpha)$ .

$(\Leftarrow)$ : Suppose  $\text{EL}(M, M, \text{ORD})$  holds.

- ▶  $M \models \varphi(a, \beta) \wedge \neg\varphi(b, \beta)$  for some ordinal  $\beta$ .
- ▶  $V_\gamma \models \varphi(a, \beta) \wedge \neg\varphi(b, \beta)$  (by reflection).
- ▶  $\delta$  codes  $\langle \beta, \gamma \rangle$ .
- ▶  $\gamma, \beta$  are definable in  $V_{\delta+1}$  without parameters.
- ▶  $a$  and  $b$  are discernible in  $V_{\delta+1}$ .  $\square$

## Leibniz-Mycielski axiom (continued)

### Corollary:

- LM is **first-order** expressible.
- LM is equivalent to **LM\***: If  $a \neq b$  in  $M \models \text{ZF}$ , then  $M$  **satisfies** that  $a$  and  $b$  are discernible in some  $V_\alpha$ .

**Observation:**  $V = \text{HOD} \rightarrow \text{LM}$  over ZF.

**Question:** Is LM **equivalent** to  $V = \text{HOD}$  over ZF(C)?

## Leibniz-Mycielski axiom (continued)

**Theorem:** (Enayat, '04)  $M \models \text{ZF} + \text{LM}$  iff  $M$  has a definable injection from  $M$  into subsets of ordinals

$$F : M \xrightarrow{1-1} < \text{ORD}_2.$$

**Proof:**

( $\Rightarrow$ ): Suppose LM holds.

- Fix a set  $a$ . Let  $\beta_a$  be least above  $\text{rank}(a)$  such that  $V_{\beta_a} \prec_{\Sigma_2} M$ .
- $T_a$  consists of pairs  $\langle \ulcorner \varphi \urcorner, \alpha \rangle$  such that  $V_\alpha \models \ulcorner \varphi(a) \urcorner$  for some  $\alpha < \beta_a$ .
- Suppose  $a \neq b$  and (wlog)  $\text{rank}(a) \leq \text{rank}(b)$ .
- $V_\gamma \models \varphi(a) \wedge \neg \varphi(b)$  for some  $\gamma < \beta_b$ .
- $\langle \ulcorner \neg \varphi \urcorner, \gamma \rangle$  is in  $T_b$  but not in  $T_a$ .
- $T_a \neq T_b$ .
- $F(a) = T_a$  (viewed as subset of ordinals via coding).

( $\Leftarrow$ ): Suppose  $F : M \xrightarrow{1-1} < \text{ORD}_2$  and  $a \neq b$ . Then (wlog)  $\alpha \in F(a)$ , but  $\alpha \notin F(b)$ .  $\square$

## LM and choice principles

**Corollary:** If  $M \models \text{ZF} + \text{LM}$ , then  $M$  has a definable linear ordering.

**Proof:**  ${}^{<\text{ORD}}2$  is linearly ordered lexicographically.  $\square$

The existence of a definable linear ordering is a **weak global choice principle**.

**Theorem:** (Easton, '64) There are **models of ZFC without a definable linear ordering**.

**Proof:** Consider the class forcing extension  $V[G]$ , where a Cohen subset is added to every regular cardinal.  $\square$ .

**Question:** Are there models of **ZFC** having a **definable linear ordering**, but **no definable well-ordering** ( $V = \text{HOD}$  fails)?

**Question:** Is **LM** equivalent over  $\text{ZF}(C)$  to the **existence of a definable linear ordering**?

## LM and choice principles (continued)

**Theorem:** (Solovay, '04) LM does not imply (even countable) AC over ZF.

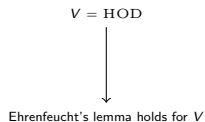
**Proof:** Similar to Cohen's argument that countable AC fails in  $L(\{c_i \mid i < \omega\})$ , where  $c_i$  are the Cohen reals explicitly added by  $\prod_{i < \omega} \mathbb{P}$ . But uses Jensen reals.

Work in  $L$ .

- Jensen forcing  $\mathbb{Q}$ :
  - ▶ A subset of Sacks forcing which has the ccc in  $L$  (constructed using  $\diamond$ ).
  - ▶ Adds a unique  $\mathbb{Q}$ -generic real over  $L$ .
  - ▶ The collection of all Jensen reals in any  $V$  is  $\Pi_2^1$ -definable.
- $\mathbb{Q}^* = \prod_{i < \omega} \mathbb{Q}$  (finite-support).
- **Theorem:** (Kanovei, '14) If  $G \subseteq \mathbb{Q}^*$  is  $L$ -generic, then the Jensen reals  $\{c_i \mid i < \omega\}$  added explicitly by  $G$  are the only Jensen reals in  $L[G]$ .
- Consider  $L(C)$ , where  $C = \{c_i \mid i < \omega\}$ .
  - ▶ countable AC fails in  $L(C)$ .
  - ▶  $C$  is definable in  $L(C)$  (as the collection of all Jensen reals).
  - ▶ Every set in  $L(C)$  is ordinal definable from a unique minimal finite subset of  $C$ .
  - ▶  $L(C)$  has a definable injection into subsets of ordinals.
  - ▶  $L(C)$  is a model of LM.  $\square$

# Questions

**Question:** Suppose  $V \models \text{ZFC}$ . Which arrows reverse?



Thank you!